

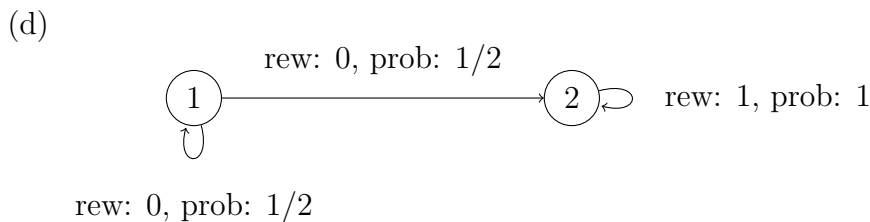
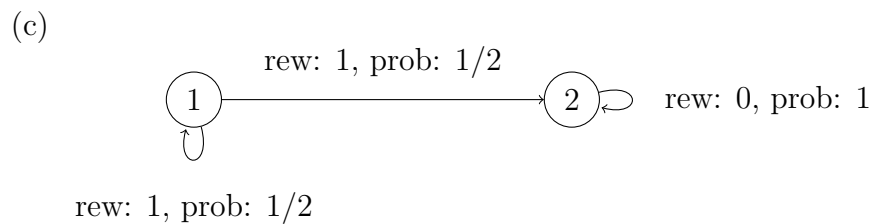
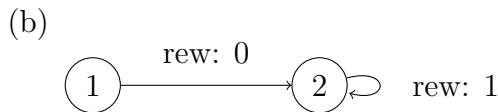
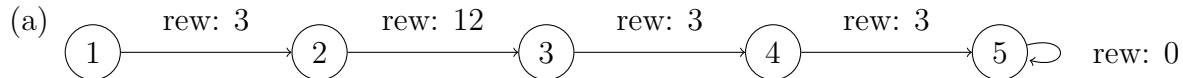
Algorithms and Uncertainty

Summer Term 2020

Exercise Set 7

Exercise 1: (3+2+2+2 Points)

For the following single-armed bandits, give the fair charges of all states. Unless states otherwise, the transitions are deterministic. Justify your statements if necessary. For part (a), consider $\gamma = \frac{1}{2}$; for the remaining parts an arbitrary $\gamma \in (0, 1)$.



Exercise 2: (4 Points)

Consider the following explore-exploit algorithm. In the first $\frac{T}{2}$ steps (so $k = \frac{T}{2n}$), we explore. Afterwards, we exploit the most promising arm. Use the approach from the lecture to derive an upper-bound for the expected regret of this algorithm.

Exercise 3: (8 Points)

Use the one-sided version of Hoeffding's inequality to show a regret bound for UCB1 of $\sum_{i \neq i^*} \frac{4 \ln T}{\Delta_i} + 2\Delta_{i^*}$. The one-sided version of Hoeffding's inequality is as follows: Let Z_1, \dots, Z_N be independent random variables such that $a_i \leq Z_i \leq b_i$ with probability 1. Let $\bar{Z} = \frac{1}{N} \sum_{i=1}^N Z_i$ be their average. Then for all $\gamma \geq 0$

$$\Pr [\bar{Z} - \mathbf{E}[\bar{Z}] \geq \gamma] \leq \exp \left(- \frac{2N^2\gamma^2}{\sum_{i=1}^N (b_i - a_i)^2} \right).$$

Hint: Note that the one-sided version of Hoeffding's inequality also implies a bound on $\Pr [\bar{Z} \leq \mathbf{E}[\bar{Z}] - \gamma]$.